arrow-kt.io



Super-charge your build with

Arrow Analysis

Alejandro Serrano Mena – Tweag I/O

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- Mark possibly-nullable values with ?
- Safe access via Elvis operators ?. and ?:

One of the most celebrated Kotlin features is **nullability analysis**

- Mark possibly-nullable values with ?
- Safe access via Elvis operators ?. and ?:
- Powerful static data and control flow analysis

if (list ≠ null) {
 list.map { it + 1 } // no ?. required
}

Why stop there?

Nullability analysis saves us from NullReferenceException

Wrong indexing leads to IndexOutOfBoundsException
 list.get(2) // what if fewer elements?

Bad initialization leads to IllegalArgumentException
Person(age = -1) // not really an age

Arrow Analysis to the rescue!

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Flow-aware static analyzer for pre-, postconditions and invariants for Kotlin



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Flow-aware static analyzer for pre-, postconditions and invariants for Kotlin

- Runs as part of your build, as a compiler plug-in
- Checks information provided as annotations
 - As with nullable types and require
- Understands the order and branching of your code
 As opposed to a simpler linter

Counting semaphore

Incrementing the semaphore should only happen for **positive** numbers

import arrow.analysis.pre

fun increment(x: Int): Int { pre(x > 0) { "value must be positive" } return x + 1

Counting semaphore

Incrementing the semaphore should only happen for **positive** numbers

val example = increment(-1)

e: pre-condition `value must be positive`
 is not satisfied in `increment(-1)`
 → unsatisfiable constraint: `(-1 > 0)`

□ Flow-awareness

Flow (if, when, other calls) is taken into account when deciding whether a condition is satisfied

- val new = if (current > 1) {
 increment(current) // fine
- } else {
 log("weird...")
 0

Ġ After the call

The following code is rejected

increment(increment(1))

because increment makes no promises about its result

fun increment(x: Int): Int {
 pre(x > 0) { "value must be positive" }
 return x + 1
}

Ġ After the call

 \mathbf{F}

The following code is rejected

increment(increment(1))



because increment makes no promises about its result

fun increment(x: Int): Int {
 pre(x > 0) { "value must be positive" }
 return x + 1



The following code is accepted

increment(increment(1))

because increment makes promises about its result

fun increment(x: Int): Int {
 pre(x > 0) { "value must be positive" }
 return (x + 1)
 .post({ it > 0 }) { "result is positive" }
}

🔄 Post-conditions are checked

fun increment(x: Int): Int { val new = if (x < 0) { **0** // fails to satisfy the post-condition } else { x + 1 } return new.post({ it > 0 }) { "positive" }

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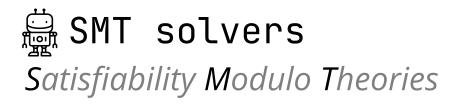


How does Arrow Analysis know that this holds?

🔅 The magic The reasoning engine

How does Arrow Analysis know that this holds?

```
fun increment(x: Int): Int {
  pre(x > 0) { "value must be positive" }
  return (x + 1)
    .post({ it > 1 }) { "result is positive" }
}
 x > 0 \& result = x + 1
   \implies result > 1
                                 SMT solver
```

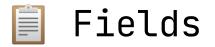


Specialized software for automatic reasoning

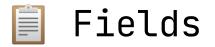
Very fast but limited to a few theories: #numbers # bitvectors * regular expressions ...

Several industrial-grade solvers **Z** (CVC5, SMTInterpol

We interface with them using java-smt



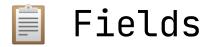
fun List<Double>.average() {
 pre(this.isNotEmpty()) { "list not empty" }
 return this.sum() / this.size
}



fun List<Double>.average() {
 pre(this.isNotEmpty()) { "list not empty" }
 return this.sum() / this.size
}

Otherwise ArithmeticException may be thrown

Odivision by zero



fun List<Double>.average() {
 pre(this.isNotEmpty()) { "list not empty" }
 return this.sum() / this.size
}

Arrow Analysis knows about the relationship

this.isNotEmpty() \iff this.size > 0



fun List<Double () { pre(this. return th }</pre>

Arrow Analysis knows about the relationship

this.isNotEmpty() \iff this.size > 0



Encode pre- and post-conditions separate from implementation

```
@Law
inline fun <E> List<E>.getLaw(index: Int): E {
  pre(index ≥ 0 && index < size) { "within bounds" }
  return get(index)
}</pre>
```

Inspired by TypeScript declaration files, but using JVM-specific features, like annotations



Encode pre- and post-conditions separate from implementation

```
@Law
inline fun <E> List<E>.getLaw(index: Int): E {
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}</pre>
```

Currently annotated: standard library
Next in our radar: (much larger, help more than welcome)



Conditions which apply to a whole type

data class Positive(val value: Int) { init { require(value > 0) } }

Useful to avoid repeating conditions again and again data class Person(val age: Positive, ...)

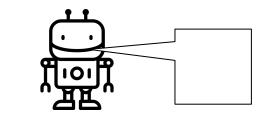


}

Invariants are assumed when using that type

```
data class Positive(val value: Int) {
    init { require(value > 0) }
```

```
operator fun plus(other: Positive) =
   Positive(this.value + other.value)
```



📑 Invariants in value classes

More compile checks with the same runtime performance!

```
@JvmInline
value class Positive(val value: Int) {
    init { require(value > 0) }
    ...
}
```

Defining your domain more strictly is A Good Thing™

🐘 Practical matters

How do I add Arrow Analysis to my project?

Our Gradle plug-in adds it to your build Instructions available at arrow-kt.io/docs/analysis

How much does it add to my compile times?

Our (very) preliminary results say that around 3x

Does it integrate with IntelliJ?

Tighter integration

Does it integrate with IntelliJ?

Interpretended in the second secon	: faile 20 sec, 944 ms
Gerrors :: CompileKotlinJvm 2 errors	3 sec, 687 ms
# example.kt src/commonMain/kc	lin 2 errors
9 pre-condition `index within bo	unds` is not satisfied
Operation `index within boom index within boo	inds` is not satisfied

The new Kotlin Frontend IR promises better integration

Our plan is to migrate once the API stabilizes

Tighter integration

Does it integrate with GitHub? 😉



0 // + numbers[0] + numbers[1] // <- problems!

0 + numbers[0] + numbers[1] // <- problems!

```
X Check failure on line 4 in src/commonMain/kotlin/example.kt
```

```
Code scanning
```

A pre-condition for a (method, property, function) is are not satisfied O Error

pre-condition index within bounds is not satisfied in numbers[0]

-> unsatisfiable constraint: true && (0 < (numbers.size))

->

4

0 bound to param index in kotlin.collections.List.get

-> in branch: 0 != null, numbers != null

Show more details



CORE

Functional companion to Kotlin's Standard Library



Functional Effects Framework

Coroutines



OPTICS

Deep access and transformations over immutable data



ANALYSIS

Pre-, post-condition, and invariant checks for your Kotlin code

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